

## MOTIVATION

- Failures in energy distribution networks (EDN) can cause enormous damage but are difficult to monitor.
- Although sensors are deployed along the pipelines for safety purposes, these sensors are vulnerable to incidents, natural events, or malicious attacks.
- Mechanisms to detect or tolerate sensor failures in EDN are desirable.

## PROBLEM DESCRIPTION

- Our goal is to design a data collection protocol in EDN that is:
  - fast and scalable;
  - resilient, in that it guarantees data availability at a sensor as long as the sensor does not fail; and
  - secure, such that it (1) prevents eavesdropping of data, and (2) ensures integrity of data.
- We consider the following scenario in our protocol:
  - An “honest-but-curious” mobile data collector (DC) collects data from pipeline sensors (PSs), and eventually delivers data to a pipeline manager (PM).
  - Sensors can fail in any stage of the protocol.

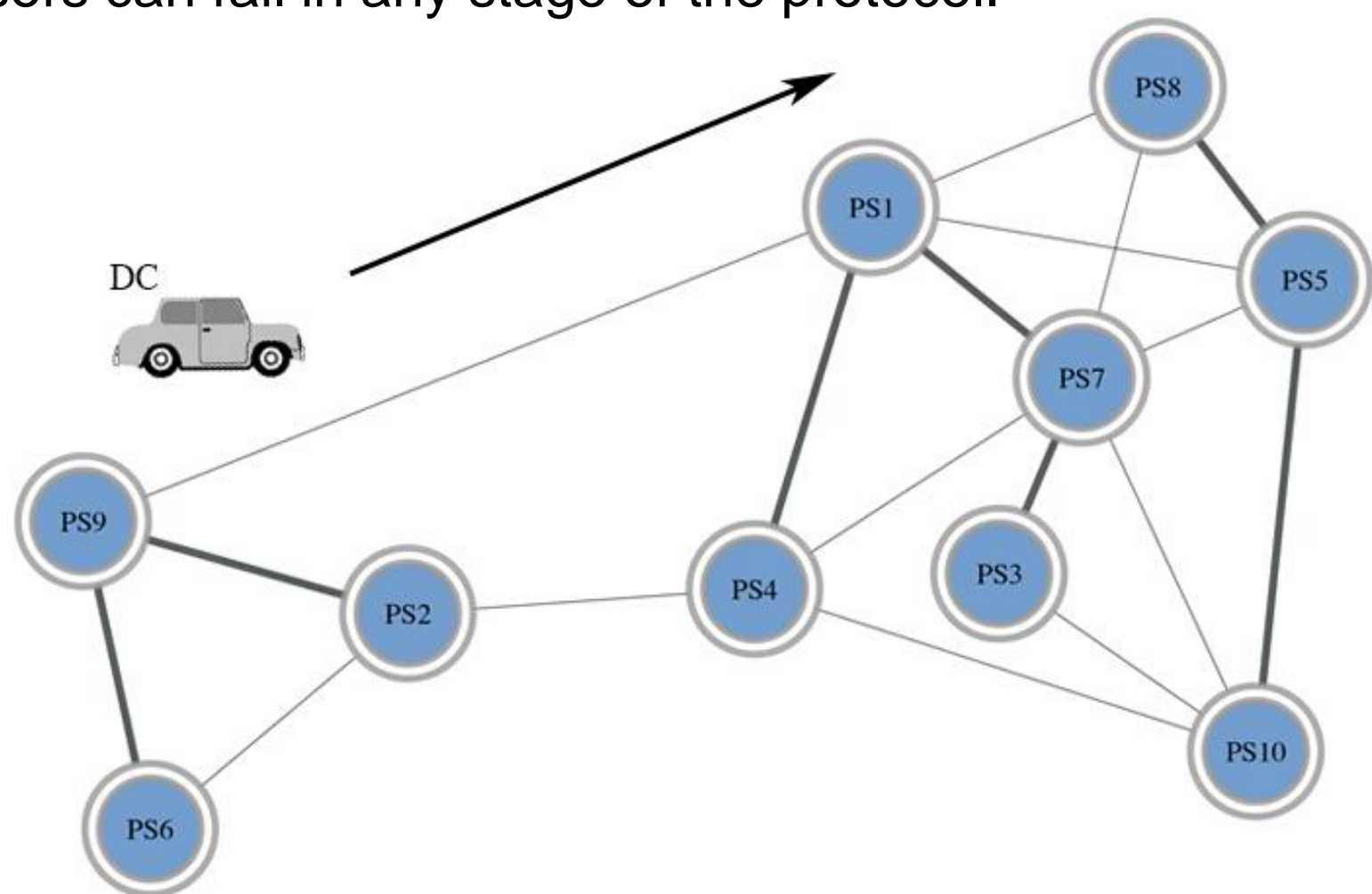


Fig. 1. Data Collection in EDN

## RELATED APPROACHES

- Secure tree-based data collection protocol in the smart grid [1].
- Scenario: in smart grid, data collector (DC) collects data from measurement devices (MD) and delivers to power operator (PO).

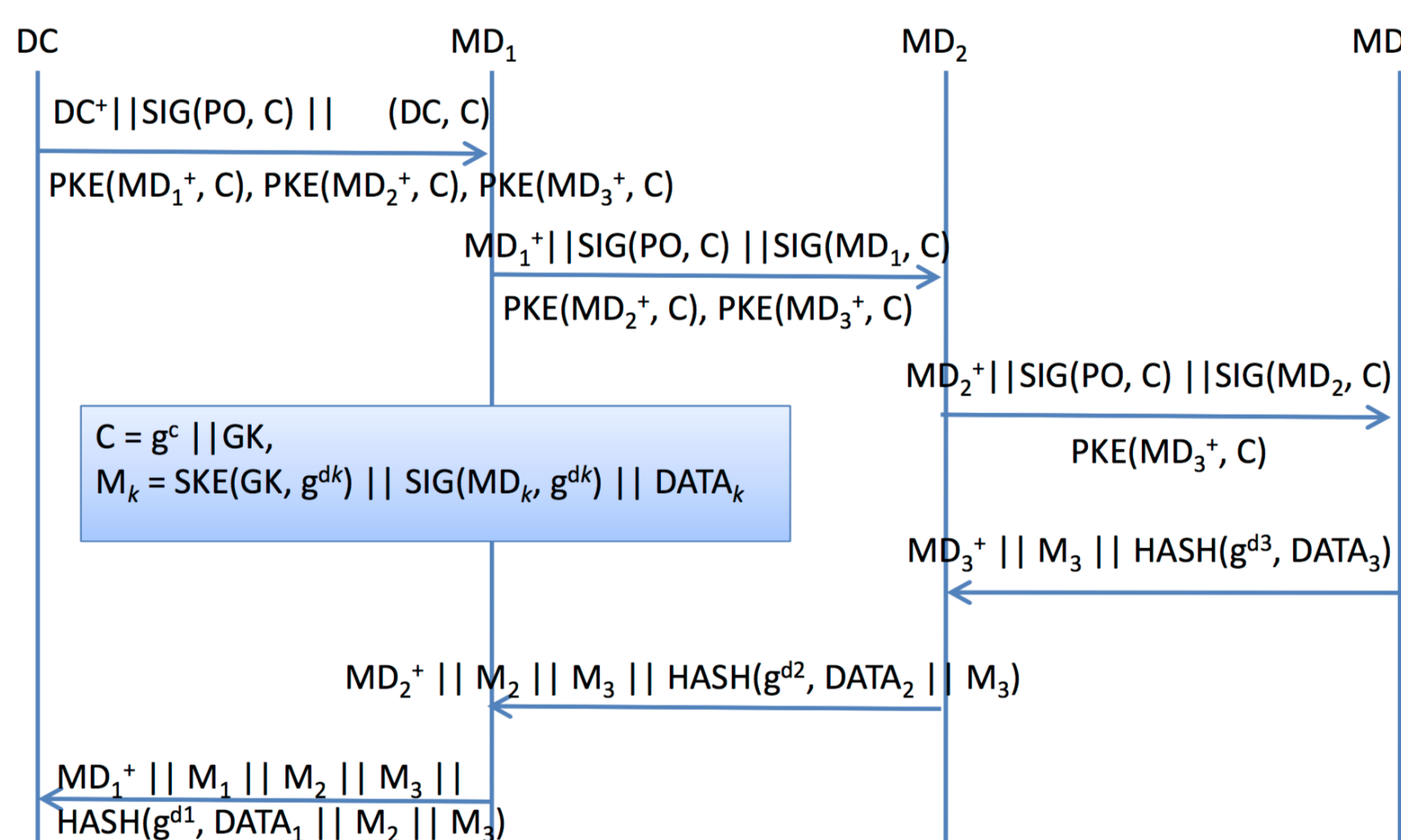


Fig. 2. Data Collection on a Tree Branch

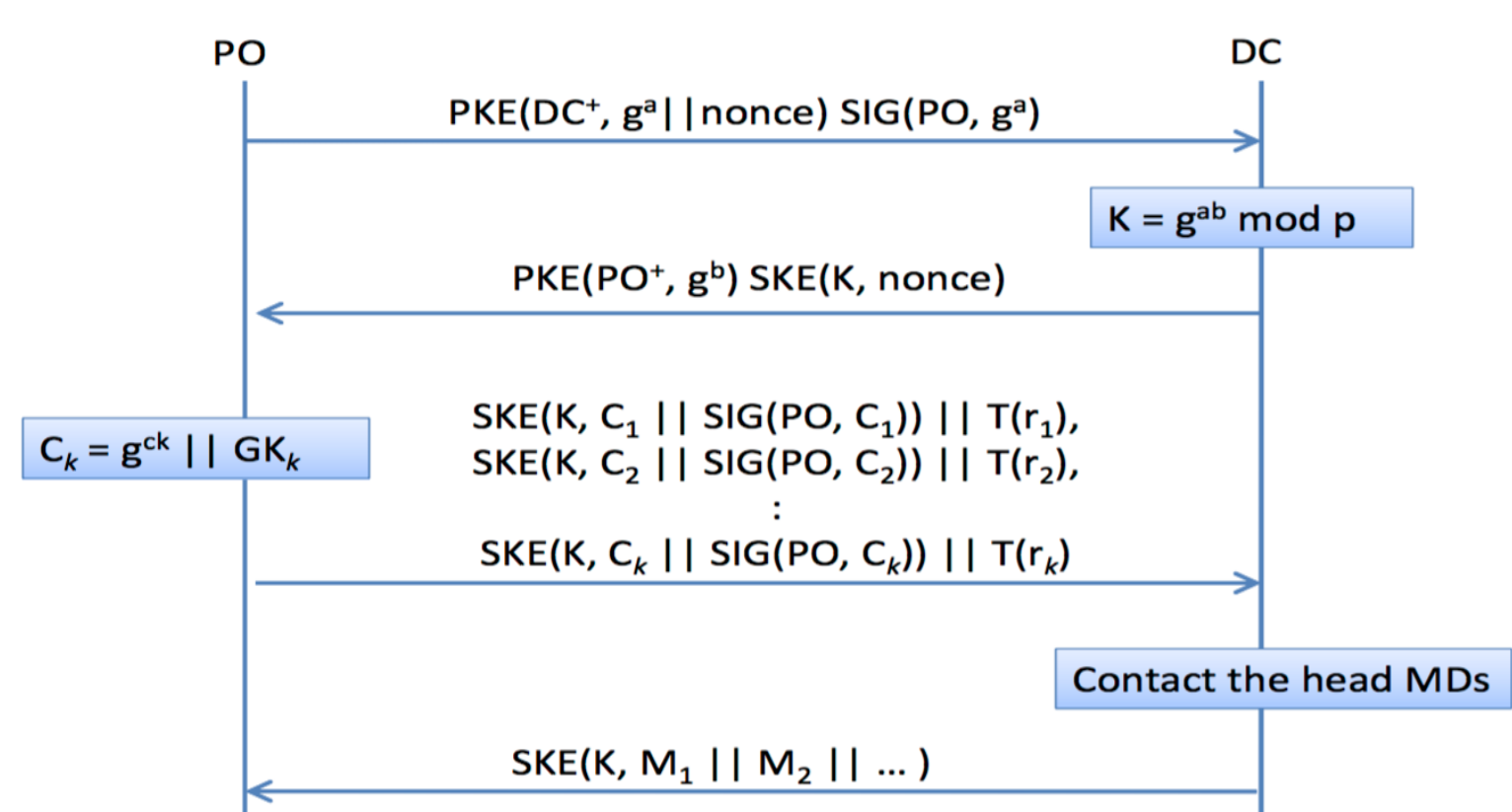


Fig. 3. Communication Between PO and DC

[1] Haiming Jin, et al. “Secure data collection in constrained tree-based smart grid environments.” *Proc. 2014 IEEE International Conference on Smart Grid Communications (SmartGridComm)*. IEEE, 2014.

## RESEARCH PLAN

- Formulate resilient data collection as a tree-based disjoint-path backup integer optimization problem, referred to as the *Resilient Tree Collection (RTC)* problem.
- Provide an algorithm to solve the optimization problem.
- Design a protocol to achieve secure and scalable data collection with/without failure.
- Evaluate the performance of the data collection scheme on real-world dataset.

## RESILIENT TREE COLLECTION (RTC)

- Notation

	Notation	Description
	$G = \langle \mathcal{R}, \mathcal{S}, \mathcal{E} \rangle$	$G$ -directed graph, $\mathcal{R}$ -set of candidate roots, $\mathcal{S}$ -set of PSs, $\mathcal{E}$ -set of links
Constants	$c_{i,j}$	Energy consumption of transmitting 1 bit through link $(i, j)$ .
	$p_{i,j}$	Probability that link $(i, j)$ fails.
	$N_{th}$	Maximum number of nodes that can share a group key.
Variables	$x_{i,j}^s$	= 1 if link $(i, j)$ is in the primary path of PS $s$ ; = 0 otherwise.
	$y_{i,j}^s$	= 1 if link $(i, j)$ is in the backup path of PS $s$ ; = 0 otherwise.

- Formulation

$$\min \sum_{j \in \mathcal{R}} \sum_{m \in \mathcal{S}} (p^s x_{i,j}^s + (1 - p^s) y_{i,j}^s) c_{i,j},$$

$$\text{s.t. } \sum_{i \in \mathcal{S}} x_{i,n}^s - \sum_{i \in \text{SUR}} x_{n,i}^s = \begin{cases} -1 & \text{if } n = s \\ 0 & \text{if } n \neq s \end{cases}, \forall s \in \mathcal{S}, \forall n \in \mathcal{S}, \quad \text{Flow}$$

$$\sum_{n \in \mathcal{R}} \sum_{i \in \mathcal{S}} x_{i,n}^s = 1, \forall s \in \mathcal{S},$$

$$x_{i,j}^s + y_{i,j}^s \leq 1, \forall s \in \mathcal{S}, \forall (i, j) \in \mathcal{E}, \quad \text{Disjoint}$$

$$x_{i,j}^s \leq x_{i,j}^i, \forall s \in \mathcal{S}, \forall (i, j) \in \mathcal{E},$$

$$\sum_{i,s \in \mathcal{S}} x_{i,j}^s + y_{i,j}^s \leq N_{th}, \forall j \in \mathcal{R} \quad \text{Security}$$

where  $p^s = \prod_{(i,j) \in \mathcal{E}} (1 - p_{i,j}^s x_{i,j}^s)$ .

## ALGORITHM

- Solve RTC (NP-hard) by rounding relaxation linear programming.

### Algorithm 1: Rounding Relaxation for RTC

**Input:**  $G, \{c_{i,j}\}, \{p_{i,j}\}, N_{th}$   
**Output:**  $\{x_{i,j}^{m*}\}, \{y_{i,j}^{m*}\}$

- while true do
- Solve LP relaxation for RTC, get optimal solution  $\{x_{i,j}^{m*}\}, \{y_{i,j}^{m*}\}$ ;
- Set  $x_{i,j}^s \leftarrow 0$  for all  $x_{i,j}^{m*} = 0$ ,  $y_{i,j}^s \leftarrow 0$  for all  $y_{i,j}^{m*} = 0$ ;
- Set  $x_{i,j}^s \leftarrow 1$  (or  $y$ ) for the largest  $x_{i,j}^{m*}$  (or  $y$ );
- if all  $\{x_{i,j}^s\}, \{y_{i,j}^s\}$  are set then
- break;
- end
- end
- return  $\{x_{i,j}^s\}, \{y_{i,j}^s\}$

## FUTURE EFFORTS

- Design protocol to achieve secure and scalable data collection under failure.
- Evaluate the performance of proposed scheme on real-world dataset.
- Consider dynamic scheduling under failure.